BUILD YOUR OWN BASIC

by Dennis Allison & Others (reprinted from *People's Computer Company* Vol. 3, No.4)

A DO IT YOURSELF KIT FOR BASIC??

Yes, available from PCC with this newspaper and a lot of your time. This is the beginning of a series of articles in which we will work our way through the design and implementation of a reasonable BASIC system for your brand X computer. We'll be working on computers based on the INTEL 8008 and 8080 microprocessors. But it doesn't make much difference — if your machine is the ZORT 9901 or ACME X you can still build a BASIC for it. But remember, it's a hard job and will take lots of time particularly if you haven't done it before. A good BASIC system could easily take one man six months!

We'd like everyone interested to participate in the design. While we could do it all ourselves, (we have done it before) your ideas may be better than ours. Maybe we can save you, or you can save us, a lot of work or problems. Write us and we'll publish your letter and comments.

WHICH BASIC?

There is not any one standard BASIC (yet). The question is which BASIC should we choose to implement. A smaller (fewer statements, fewer features) BASIC is easier to implement and (more important) takes less space in the computer. Memory is still expensive so the smaller the better. Yet maybe we can't give up some goodies like string variables, dynamic array allocation, and so on.

There is a standard version of BASIC which is to be the minimal language which can be called BASIC. It's a pretty big language with lots of goodies. Maybe too big. Is there any advantage to being compatible with, say, the EDU BASICS? We don't have to make any decision yet; but the time will come ...

COMPILER OR INTERPRETER?

We favor using an interpreter. An interpreter is a program which will execute the BASIC program from its textual representation. The program you write is the one which gets executed. A compiler converts the BASIC program into the machine code for the machine it is to run on. Compiled code is a lot faster, but requires more space and some kind of mass storage device (tape or disk). Interpretative BASIC is the most common on small machines.

HOW MUCH MEMORY? AND ... WHAT KIND?

Can we make some guesses about how big the BASIC system will be? Only if you don't hold us to it. Suppose we want to be able to run a 50 line BASIC program. We need about 800 bytes to store the program, another 60 or so bytes for storing program values (all numeric) without leaving any space for the interpreter and its special data. Past experience has shown that something like 6 to 8 Kbytes are needed for a minimum BASIC interpreter and that at least 12K bytes are needed for a comfortable system. That's a lot of memory, but not too much more than you need to run the assembler. A lot of BASIC could be put into ROM (Read Only Memory) once developed and checked out. ROM is a lot cheaper than RAM (Read and Write) memory, but you can't change it. It's lots better to make sure everything works first.

DIRECT MODE?

Some kind of "desk calculator" mode of opera-

tion would be nice. At least, we would like to

This feature makes it easier to find and gently

terminate the existence of "bugs."

be able to look at and set different variables in a

program and restart execution at any given point.

But... if we can agree on some chunks of code and get it properly checked out, some enterprising person out there might make a few thousand ROMs and save us all some \$\$\$. Let's see now... how about ROMs for floating point arithmetic, integer arithmetic, Teletype I/O...

DATA STRUCTURES

Data structures are places to put things so you can find them or use them later. BASIC has at least three important ones: a symbol table which looks up a program name, A or Z9 or A\$, with its value. If we had a big computer where space was not a huge problem, we could simply preallocate all storage since BASIC provides for only 312 different names excluding arrays. When memory is so costly this doesn't make much sense. Somewhere, also, we've got to store the names which BASIC is going to need to know, names like LET and GO TO and IF. This table gets pretty big when there are lots of statements.

Lastly, we need some information about what is a legal BASIC statement and which error to report when it isn't. These tables are called parsing tables since they control the decomposition of the program into its component parts.

STRATEGY

Divide and Conquer is the programmers maxim. BASIC will consist of a lot of smaller pieces which communicate with each other. These pieces themselves consist of smaller pieces which themselves consist of smaller pieces, and so forth down to the actual code. A large problem is made manageable by cutting it into pieces.

What are the pieces, the building blocks of BASIC? We see a bunch of them:

- * a supervisor which determines what is to be done next. It receives control when BASIC is loaded.
- * a program and line editor. This program collects lines as they are entered from the keyboard and puts them into a part of computer memory for later use.
- *a line executor routine which executes a single BASIC statement, whatever that is
- *a line sequence which determines which line is to be executed next.
- *a floating point package to provide floating point on a machine without the hardware.
- *terminal I/O handler to input and output information from the Teletype and provide simple editing (backspace and line deletion).
- *a function package to provide all the BASIC functions (RND, INT, TAB, etc.) *an error handling routine (part of the supervisor).
- *a memory management program which provides dynamic allocation data objects.

These are the major ones. As we get futher into the system we'll begin to see others and we'll begin to be able to more fully define the function of each of these modules.

TINY BASIC

Pretend you are 7 years old and don't care much about floating point arithmetic (what's that?), logarithms, sines, matrix inversion, nuclear reactor calculations and stuff like that.

And ... your home computer is kinda small, not too much memory. Maybe its a MARK-8 or an ALTAIR 8800 with less than 4K bytes and a TV typewriter for input and output.

You would like to use it for homework, math recreations and games like NUMBER, STARS, TRAP, HURKLE, SNARK, BAGELS, ...

Consider then, TINY BASIC

- Integer arithmetic only 8 bits? 16 bits?
- 26 variables: A, B, C, D, ..., Z
- The RND function of course!
- Seven BASIC statement types INPUT

PRINT LET

GO TO

- IF
- GOSUB
- RETURN
- Strings? OK in PRINT statements, not OK otherwise.

BUILD YOUR OWN BASIC--REVIVED

(reprinted from People's Computer Company Vol. 4, No. 1)

WHAT IS TINY BASIC???

TINY BASIC is a very simplified form of BASIC which can be implemented easily on a microcomputer. Some of its features are:

Integer arithmetic 16 bits only

26 variables (A, B, ..., Z)

Seven BASIC statements



INPUT PRINT LET GOTO IF GOSUB RETURN

Strings only in PRINT statements

Only 256 line programs (if you've got that much memory)

Only a few functions including RND

It's not really BASIC but it looks and acts a lot like it. I'll be good to play with on your ALTAIR or whatever; better, you can change it to match your requirements and needs.

TINY BASIC LIVES!!!

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wh the We are working on a version of TINY BASIC to run on the INTEL 8080. It will be an interpretive system designed to be as conservative of memory as possible. The interpreter will be programmed in assembly language, but we'll try to provide adequate descriptions of our intent to allow the same system to be programmed for most any other machine. The next issue of PCC will devote a number of pages to this project.

In the meantime, read one of these. Compiler Construction For Digital Computers, David Gries, Wiley, 1971 493 pages, \$14.95

Theory & Application of a Bottom-Up Syntax Directed Translator Harvey Abramson, Academic Press, 1973, 160 pages, \$11.00

Compiling Techniques, F.R.A. Hopgood, American Elsevier, 126 pages \$6.50

A BASIC Language Interpreter for the Intel 8008 Microprocessor A.C. Weaver, M.H. Tindall, R.L. Danielson. University of Illinois Computer Science Dept, Urbana IL 61801. June 1974. Report No. UIUCDCS-R-74-658. Distributed by National Technical Information Service, U.S. Commerce Dept, Springfield VA 22151. \$4.25.

A BASIC language interpreter has been designed for use in a microprocessor environment. This report discusses the development of 1) an elaborate text editor and 2) a table-driven interpreter. The entire system, including text editor, interpreter, user text buffer, and full floating point arithmetic routines fits in 16K 8-bit words.

The TINY BASIC proposal for small home computers was of great interest to me. The lack of floating point arithmetic however, tends to limit its usefulness for my objectives.

As a matter of a suggestion, consideration should be given to the optional inclusion of floating point arithmetic, logarithm and trigonmetric calculation capability via a scientific calculator chip interface.[†]

The inclusion of such an option would tend to extend

the interpreter to users who desire these complex calculation capabilities. A number of calculator chip proposals have been made, with the Suding unit being of the most interest.

Thank you for the note of 13 June, regarding my letter on the Tiny BASIC article (PCC Vol. 3 No. 4). It was with regret that I learned that the series was not continued in the next volume. Even though few responded to the article published, conceptually the knowledge and principles which would be disseminated regarding a limited lexicon, high level programming language are of importance to the *independent* avocational microcomputer community.

At this time, PCC may not have a wide distribution in the avocation microcomputer community. This could be possibly the cause for the low number of respondies. Never the less, this should not detract from the dissemination and importance of concepts and principles which are of significance.

The thrust of my letter of 15 April, 1975, was to suggest a mechanism for the inclusion of F.P. in a limited lexicon and memory consumptive BASIC. I hope that the implication that F.P. must be included was not read into my letter.

It is my interest that information, concepts and the principles of compiler/interpreter construction as it related to microcomputers be available to the limited budget avocational user. The MITS BASIC, which you brought up, appears from my viewpoint to be a *licensed*, blackbox program which is not currently available to: (a) 8008 users, (b) IMP-16 users, (c) independent 8080 users (except at a very large expense) or (d) MC6800 users who will shortly be on linc.

Presently it appears that microcomputer compiler interpretor function langauges will be coming available from a number of sources (MITS, NITS, Processor Technology and etc.). However, few will probably deal in the conceptualizations which are the basis of the interpreter. Information which will fill the void in the interpreter construction knowledge held by the avocation builder, should be made available.

I strongly urge that the series started with Vol. 3 No. 4 article be continued. Possibly the hardware, peripheral, machine programming difficulties incurred by the microcomputer builder, is prohibiting a major contribution at this time. However, I would expect that by Autumn a number of builders should have their construction and peripheral difficulties far enough along to start thinking about higher level languages. The previous objective for the article series sounds reasonable. It was not my purpose in submitting the letter to detract from the objective of a very limited lexicon BASIC, ie., to be attractive and usable by the young and beginner due to its simplicity. If wives, children, neighbors or anyone who is not machine language or programming oriented is expected to use a home-base unit created under a restrained hudget a high level language will be a necessity. It is with this foresight that I encourage the continuance of the "Build Your Own BASIC" series. This issue aside, I would like to encourage the PCC to continue the quite creditable activities which have been its order of business with regard to avocational Michael Christoffer computing. 4139 12th NE No. 400

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† Please see Dr Robert Suding's article on p. 18

DESIGN NOTES FOR TINY BASIC

by Dennis Allison, happy Lady, & friends (reprinted from *People's Computer Company* Vol. 4, No. 2)

SOME MOTIVATIONS

A lot of people have just gotten into having their own computer. Often they don't know too much about software and particularly systems software, but would like to be able to program in something other than machine language. The TINY BASIC project is aimed at you if you are one of these people. Our goals are very limited to provide a minimal BASIC-like language for writing simple programs. Later we may make it more complicated, but now the name of the game is **keep it simple**. That translates to a limited language (no floating point, no sines and cosines, no arrays, etc.) and even this is a pretty difficult undertaking.

Originally we had planned to limit ourselves to the 8080, but with a variety of new machines appearing at very low prices, we have decided to try to make a portable TINY 3ASIC system even at the cost of some efficiency. Most of the language processor will be written in a pseudo language which is good for writing interpreters like TINY BASIC. This pseudo language (which interprets TINY BASIC) will then itself be implemented interpretively. To implement TINY BASIC on a new machine, one simply writes a simple interpreter for this pseudo language and not a whole interpreter for TINY BASIC.

We'd like this to be a participatory design project. This sequence of design notes follows the project which we are doing here at PCC. There may well be errors in content and concept. If you're making a BASIC along with us, we'd appreciate your help and your corrections.

Incidentally, were we building a production interpreter or compiler, we would probably structure the whole system quite differently. We chose this scheme because it is easy for people to change without access to specialized tools like parser generator programs.

THE TINY BASIC LANGUAGE

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There isn't much to it. TINY BASIC looks like BASIC but all variables are integers There are no functions yet (we plan to add RND, TAB, and some others later). Statement numbers must be between 1 and 255 so we can store them in a single byte. LIST only works on the whole program. There is no FOR-NEXT statement. We've tried to simplify the language to the point where it will fit into a very small memory so impecunious. tyros can use the system.

The boxes shown define the language. The guide gives a quick reference to what we will include. The formal grammar defines exactly what is a legal TINY BASIC statement. The grammar is important because our interpreter design will be based upon it.

IT'S ALL DONE WITH MIRRORS-----OR HOW TINY BASIC WORKS

All the variables in TINY BASIC: the control information as to which statement is presently being executed and how the next statement is to be found, the return addresses of active GOSUBS-----all this information constitutes the state of the TINY BASIC interpreter.

There are several procedures which act upon this state. One procedure knows how to execute any TINY BASIC statement. Given the starting point in memory of a TINY BASIC statement, it will execute it changing the state of the machine as required. For example,

100 LET S = A+6 (C) would change the value of S to the sum of the contents of the variable A and the interger 6, and sets the next line counter to whatever line follows 100, if the line exists.

A second procedure really controls the interpretation process by telling the line interpreter what to do. When TINY BASIC is loaded, this control routine performs some initialization, and then attempts to read a line of information from the console. The characters typed in are saved in a buffer, LBUF. It first checks to see if there is a leading line number. If there is, it incorporates the line into the program by first deleting the line with the same line number (if it is present) then inserting the new line if it is of nonzero length. If there is no line number present, it attempts to execute the line directly. With this strategy, all possible commands, even LIST and CLEAR and RUN are possible inside programs. Suicidal' programs are also certainly possible.

TINY BASIC GRAMMAR

The things in **bold face** stand for themselves. The names in lower case represent classes of things. '::=' is read 'is defined as'. The asterisk denotes zero or more occurances of the object to its immediate left. Parenthesis group objects. ϵ is the empty set. | denotes the alternative (the exclusive-or).

line::= number statement 🞯 statement 🕝						
statement::= PRINT expr-list						
IF expression relop expression THEN statement						
GOTO expression						
INPUT var-list						
LET var = expression						
GOSUB expression						
RETURN						
CLEAR						
LIST						
RUN						
END						
expr-list::= (string expression) (, (string expression) *)						
var-list::= var (, var)*						
expression::= $(+ - \varepsilon)$ term $((+ -)$ term)*						
term::= factor ((* /) factor)*						
factor::= var number (expression)						
var::= A B C Y Z						
number::= digit digit*						
diait::= 0] 1 [2] 8 9						
relop::=<(> = E) >(< = E) =						
A BREAK from the console will interrupt execution of the program.						
A BUEAR HOW WE CONSIGN WITH MEN OF COUNCIL OF THE PROGRAM						

IMPLEMENTATION STRATGLES AND ONIONS

When you write a program in TINY BASIC there is an abstract machine which is necessary to execute it. If you had a compiler it would make in the machine language of your computer a program which emulates that abstract machine for your program. An interpreter implements the abstract machine for the entire language and rather than translating the program once to machine code it translates it dynamically as needed. Interpreters are programs and as such have their's as abstract machines. One can find a better instruction set than that of any general purpose computer for writing a particular interpreter. Then one can write an interpreter to interpret the instructions of the interpreter which is interpreting the TINY BASIC program. And if your machine is microprogrammed (like PACE), the machine which is interpreting the interpreter interpreting the interpreter interpreting BASIC is in fact interpreted.

This multilayered, onion-like approach gains two things: the interpreter for the interpreter is smaller and simpler to write than an interpreter for all of TINY BASIC, so the resultant system is fairly portable. Secondly, since the major part of the TINY BASIC is programmed in a highly memory efficient, tailored instruction set, the interpreted TINY BASIC will be smaller than direct coding would allow. The cost is in execution speed, but there is not such a thing as a free lunch.



interpreter for TINY BASIC

interpreter for TINY BASIC's interpreter

LINE STORAGE

The TINY BASIC program is stored, except for line numbers, just as it is entered from the console. In some BASIC interpreters, the program is translated into an intermediate form which speeds execution and saves space. In the TINY BASIC environment, the code necessary to provide the

QUICK REFERENCE GUIDE FOR TINY BASIC

LINE FORMAT AND EDITING

- Lines without numbers executed immediately.
- Lines with numbers appended to program
- · Line numbers must be 1 to 25
- · Line number alone (empty line) deletes line
- Blanks are not significant, but key words must contain no unneeded blanks
- '---' deletes last character
- X^C deletes the entire line

EXECUTION CONTROL

CLEAR delete all lines and data RUN run program LIST list program

EXPRESSIONS

Operators

Arithmetic	Relational			
+	>	>=		
* /	4	4=		
	2	<>,><		
		•		

Variables A.....Z (26 only)

All arithmetic is modulo 2^{15} (± 32762)

INPUT / OUTPUT

PRINT X,Y,Z PRINT 'A STRING' PRINT 'THE ANSWER IS' INPUT X INPUT X,Y,Z

ASSIGNMENT STATEMENTS

LET X=3 LET X= -3+5*Y

CONTROL STATEMENTS

GOTO X+10 GOTO 35 GOSUB X+35 GOSUB 50 RETURN IF X > Y THEN GOTO 30

transformation would easily exceed the space saved.

When a line is read in from the console device, it is saved in a 72-byte array called LBUF (Line BUFfer). At the same time, a pointer, CP, is maintained to indicate the next available space in LBUF. Indexing is, of course, from zero.

Delete the leading blanks. If the string matches the BASIC line, advance the cursor over the matched string and execute the next IL instruction. If the match fails, continue at the IL instruction labeled lbl.

The TINY BASIC program is stored as an array called PGM in order of increasing line numbers. A pointer, PGP, indicates the first free place in the array. PGP=0 indicates an empty program; PGP must be less than the dimension of the array PGM. The PGM array must be reorganized when new lines are added, lines replaced, or lines are deleted.

Insertion and deletion are carried on simultaneously. When a new line is to be entered, the PGM array searches for a line with a line number greater than or equal to that of the new line. Notice that lines begin at PGM (0) and at PGM (j+1) for every j such that PGM (j)=[carriage return]. If the line numbers are equal, then the length of the existing line is computed. A space equal to the length of the new line is created by moving all lines with line numbers greater than that of the line being inserted up or down as appropriate. The empty line is handled as a special case in that no insertion is made.

TINY BASIC AS STORED IN MEMORY

byte in memory treated as an integer

N *

WAEND C

N N

byte treated as a character

POW

UTLN Q'20 P

N * N * N @ '36 I

ØUTHENUGOTOU1

Two different things are going on at the same time. The routines must determine if the TINY BASIC line is a legal one and determine its form according to the grammar; secondly, it must call appropriate action routines to execute the line. Consider the TINY BASIC statement:

GOTO 100 At the start of the line, the interpreter looks for BASIC key words (LET, GO, IF, RETURN, etc.) In this case, it finds GO, and then finds TO. By this time it knows that it has found a GOTO statement. It then calls the routine EXPR to obtain the destination line number of the GOTO. The expression routine calls a whole bunch of other routines, eventually leaving the number 100 (the value of the expression) in a special place, the top of the arithmetic expression stack. Since everything is legal, the XFER operator is invoked to arrange for the execution of line 100 (if it exists) as the next line to be executed.

Each TINY BASIC statement is handled similarly. Some procedural section of an IL program corresponds to tests for the statement structure and acts to execute the statement.

ENCODING

There are a number of different considerations in the TINY BASIC design which fall in this general category. The problem is to make efficient use of the bits available to store information without loosing out by requiring a too complex decoding scheme

In a number of places we have to indicate the end of a string of characters (or else we have to provide for its length somewhere). Commonly, one uses a special character (NUL = 00H for example) to indicate the end. This costs one byte per string but is easy to check. A better way depends upon the fact that ASCII code does not use the high order bit; normally it is used for parity

ONE POTENTIAL IL ENCODING



							Box 310, Menio Park CA 94029 Tage 8
on transmission. We can use it to indicate the end			A STATEMENT EXECUTOR WRITTEN IN IL				
(that is, last character) of a string. When we process the characters we must AND the character with			This program in U, will execute a TNY BASIC tatament. The operators TST, TSTV, TSTV, and PRS all one a cursor to find characteristics of the TINY BASIC CONTRACT STATEMENT OF CONTRACT STATEMENT Descent Contract Statement of the Statement Statement Statement Statement Statement Statement Statement Statement Descent Statement S				
07FH to scrub off the flag bit.			BASIC fine, the data set in the INV BASIC fine, the operation (INV), KER an ever the cursor so it points to a mother FIRY BASIC fine.				
					H. SECTION		
The interpreter opcodes can be encoded into a single byte. Operations fall into two distinct						INITIAL	125
	•	nachine language sub-	co:	NLIN GET TSTE	LINE XEC	WRITEC	R/LF
routines, and those which either call or transfer			STMT:	KVSR JMP XINI	T	; INSERT	R LINE NUWRER IT IMAY BE DELETEI
within the IL language itself. The diagram indi-						INITIAL	IZE FOR EXECUTION
cates one encoding scheme. The CALL operations					ECUTOR	IS STAT	
•		the IL instruction set.	STMT:	1.37 TSEV CALL	EXPA	PLACE P	EMENTA LETY ACE VAR ADDRESS ON AESTK. YAPR VALUE DIT AESTK.
	Addressing is shown to be relative to PC for IL operations. Given the current IL program size,			DON STOI NXT	E RE	STORE	ERROR #For NOT NEXT,
this seems adequate. If it is not, the address		S1:	TST TST	\$3, '60' \$2, 'TO'	; GOTO O	QUENCE TO NEXT. R GOSUBY O ORSub.	
		array with the ML		CAL1 DON XPEE	S2.'TO' EXPR	GETLA	BEL.
class instructi	ons.		\$2:	TST CAL	S14, 'SUB' - EXPA	GET DES	N G ROL ACAL AND JUMP. IF NO MATCH. STINATION.
	TINY BASIC IN	TERPRETIVE OPERATIONS		DON SAV XECE		SAVE RE	IF & NOT NEXT. ETURN LINE.
	TST Ibl. 'string	defete leading blanks If string matches the BASIC line, advance cursor over the	53 54:	75T 751	SB PRINT S7	TLST FO	IA QUOTE,
		matched string and execute the next IL instruction. If a	55:	PRS TST SPC	56,','	FRINT S	TRING. E MORE?
		match fails, execute the H_ instruction at the labeled Ibl.	S6:	UMP DONI		VIS HA	0 NEXT ZONE, 39 BACK. 108 HE NG 60.
i	CALL IM	Execute the IL subroutine starting at Ibl. Save the IL ad- dress fullowing the CALL on the control stack.	\$7:	NLIN NXT	E		
	RTN	Return to the IL location specified by the top of the con-		CALL PRM JAP	55	, P6)RT 11 15 THE 0	F MORE?
		trof stack.	\$8:	TST CALL	59, 'IF' E XPO	GETEXP	EMENT, PRESSION
	DONE	Report a syntax error if after detetion leading blacks the cursor is not positioned to read a carruge return,		CALL CALL CALP	, EXPH	GET EXP	INTE ORAND PUT ON STK. PRESSION, % COMPARISON-PERFORMS NEXT IF FALSE,
	JMP Ibl	Continue execution of IL at the labet specified.	59: 510:	かけ 151	STMT S17, 'INPUT"	GET NLA	TATENT, AND CALLER OTHER TEAL OF TALSE.
1	PRS	Print characters from the BASIC text up to but not including the	S10:	CALL INNU STOR	. VAR	STORE I	UMBEN FROM TTY TO AESTK." T.
	. 115	error. Governments from the BASSC text up to but not intervining the closing quote mark. If a cr is found in the program text, report an error. Move the cursor to the point following the closing quote.		75T JMP	S11,5" S10	; IS THEN ; YES,	E MORE?
	PRN		531: 512:	DONE NXF IST	S13, 'RETUR	MUST BE	G. Ce to Next. Statement,
	ะกเพ	Print number obtained by popping the top of the expres- sion stack.		DONE	Ę.	: MUST BE : RESTOR	er. E LINE NUMBER OF CALL!
	SPC	Insert spaces to move the print head to next zone.	S13:	NXT TST FIN	\$14, 'END'	: SEQUEN	CE TO NEXT STATEMENT.
	NLINE	Output CHLF to Printer,	514:	TST DONE	515, 'LIST'	; LIST CON	MMAND.
	NXT	If the present mode is direct (line number zero), then	\$15:	LST NXT IST	\$16, "RUN"	, AUN COM	MAND .
		return to line collection. Otherwise, select the next sequential line and begin interpretation.		DONE NXT	Ę		
	XFER	Test value in the top of the AE stick to be within range.	S18:	DONE DONE	S17, 'CLEAR'	;CLEAR C	OMMAND.
		If not, report an error. If so, attempt to position cursor at that line. If it exists, hegin interpretation there; if not	\$17:	ERA		;SYNTAX	EAROA.
		Réport an eiror.	EXPA;	TST CALL	EO, '' TERM		
	SAV	Place present line number on SURSTK. Report overflow as error.		NEG JMP	£1	; TEST FD ; GET VAL ; NEGATE	R UNARY .ue. It.
	RSTR	Replace current line number with value on SBRSTK. If	EO; E1;	157 CALL 151	E1, ** TERM E2, **	100000	OR MORE, IR UNARY 4.
		stack is empty, report error.	L 1.	CALL		SUM TEP	a jenna. RM,
	CMPR	Compare AESTK(SP), the top of the stack, with AESTK(SP-2) as per the relation indicated by AESTK(SP-3),	£2:	JAIP TST CALL	E1 E3,'-' TERM	ANY MO	NE? LACE TERM.
		Delete all from stock. If conclusion specifiest did not match, then perform NXT action.		SUB AlP	E3		
	INNUM	Read a number from the terminal and push its value onto	£3: T2:	RTN		; ANY MO	RE7
		the AESTK,	TERN TO:	1:	CALL TST	FACT	
	FIN	Return to the fine collect routine.	.0.		CALL	FACT	PRODUCT FACTOR.
1	ERR	Report syntax error and return to line collect routine.			MPY JMP	то	
	ADD	Replace top two elements of AESTK by their sum,	T1:		TST	T2, '/'	ANY MORE?
	SUB	Replace top two elements of AESTK by their difference.			CALL DIV	FACT	QUOTIENT FACTOR.
	NEG	Replace top of AESTK with its negative.			JMP	то	
	MUL	Replace top two elements of AESTK by their product.	E	r.	TOTA	50	
	DIV	Replace top two elements of AESTK by their quotient.	FACT	:	TSTV IND	FO	;VARIABLE. ;YES, GET THE VALUE.{
	STORE	Place the value at the top of the AESTK into the variable			RTN	C 1	· · · ·
		designated by the index specified by the value immediately below it. Delete both from the stock.	F0:		TSTN RTN	F1	';NUMBER, GET ITS VALUE.
	TSTV Ibi	Test for variable (i.e. letter) if present., Place its index	F1:		TST	F2, '('	;PARENTHESIZED EXPR.
		value onto the AESTK and continue execution at next suggested location. Otherwise, continue at lbf,			CALL TST	EXPR F2, ')'	; MATCHING PARENTHESIS.
	TSTN Ibl	Test for number, if present, place its value onto the	_		RTN	· -, ,	
		AESTK and continue execution at next suggested location, Otherwise, continue at Ibl,	F2:		ERR		; ERROR.
	IND	Replace top of stack by variable value if indexes.		101	TOT	RO, '='	
	LST	list the contents of the program area.	RELC	JF (tst Lit	но, ≃ 0	<u>j</u> m
	INIT	Performs global initialization			RTN	DA 1/1	
		Clears program area, emptys GOSUB stack, etc.	RO:		TST TST	R4, ' (' R1, '='	
	GETLINE	Input a line to LBUF.			LIT	2	;ζ= R4: TST S17, '>'
	TSTL (b)	After editing learling blacks, look for a line number. Report error of invalid; transfer to lot if not present.	R1:		RTN TST	R3, ')'	TST R5≐′ LIT 5 >>=
	INSRT	Insert line after deleting any line with same line			LÍT	3	RTN "
		number.	R3:		RTN LIT	1	אד: TST R6, '∠' ג ⊔IT 3 ג ג ג ג ג
	XINIT	Perform initialization for each stated execution. Empties AEXP stack,			RTN		R6: LIT4 ; >
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